Transactional Tasks Futures in Software Transactions

Janwillem Swalens





Overview

- 1. Crash course: futures
- 2. Crash course: transactions
- 3. Problem: nesting futures in transactions
- 4. Solution: transactional tasks

Crash course: Futures for parallelism

```
(future e) returns f
(deref f) returns result of e
```

```
(defn fib [n]
  (if (< n 2)
        n
        (let [a (fib (- n 1))
             b (fib (- n 2))]
        (+ a b))))</pre>
```

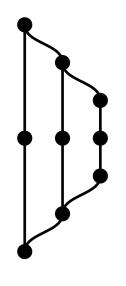
Crash course: Futures for parallelism

```
(future e) returns f
(deref f) returns result of e
```

```
Property:
```

semantically transparent

```
(future e) → e
(deref f) → f
is equivalent
(assuming no side-effects in e)
```



Overview

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- 2. Crash course: transactions
- 3. Problem: nesting futures in transactions
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Crash course: transactions for shared memory

```
Property:
(ref v)
                               serializability
(dosync e)
(deref r)
(ref-set r v)
(def checking (ref 100))
(def savings (ref 500))
(future
 (dosync
    (ref-set checking (- (deref checking) 10))
    (ref-set savings (+ (deref savings) 10))))
(future
 (dosync
    (println "You own €" (+ (deref checking)
                            (deref savings)))))
```



ref	v0	
checking	100	
savings	500	

0		
ref	v0	
checking	100	
savings	500	

T1: 0

ref	val	write?

O			T1: 0)		T2: 0	١	
ref	v0		ref	val	write?	ref	val	write?
checking	100							
savings	500							

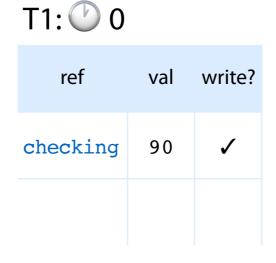
0		
ref	v0	
checking	100	
savings	500	

ref	val	write?
checking	100	

T1: 0

ref	val	write?

O		
ref	v0	
checking	100	
savings	500	



ref	val	write?

O		
ref	v0	
checking	100	
savings	500	



ref	val	write?

0
•

ref	v0	
checking	100	
savings	500	

T1: 0

ref	val	write?
checking	90	1
savings	510	1

T2: 0

ref	val	write?

0
U

ref	v0	
checking	100	
savings	500	

T1: 0

ref	val	write?
checking	90	1
savings	510	1

T2: 0

ref	val	write?
checking	100	
savings	500	

commit T1:

1. copy written values (if no more recent values exist)

0			
ref	v0	v1	
checking	100	90	
savings	500	510	

ref	val	write?
checking	90	1
savings	510	1

T1: 0

ref	val	write?
checking	100	
savings	500	

commit T1:

(7) 1

- 1. copy written values (if no more recent values exist)
- 2. increase global write-point

_			
ref	v0	v1	
checking	100	90	
savings	500	510	

ref	val	write?
checking	90	1
savings	510	1

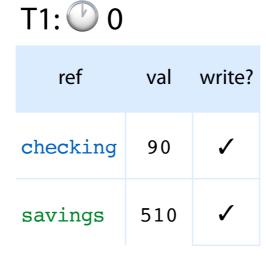
T1: 0

ref	val	write?
checking	100	
savings	500	

T1 committed

1

-			
ref	v0	v1	
checking	100	90	
savings	500	510	



ref	val	write?
checking	100	
savings	500	

commit T2:

2

- 1. copy written values (if no more recent values exist)
- 2. increase global write-point

_			
ref	v0	v1	
checking	100	90	
savings	500	510	

ref	val	write?
checking	90	1
savings	510	1

T1: 0

ref	val	write?
checking	100	
savings	500	

T2 committed

2				T1: 0			T2: 0		
ref	v0	v1		ref	val	write?	ref	val	١
checking	100	90		checking	90	1	checking	100	
savings	500	510		savings	510	1	savings	500	

Result is **serializable**. Serialization = T2; T1

```
(def checking (ref 100))
(def savings (ref 500))
(future
  (dosync
    (ref-set checking (- (deref checking) 10))
    (ref-set savings (+ (deref savings) 10))))
(future
  (dosync
    (println "You own €" (+ (deref checking)
                             (deref savings)))
    (ref-set checking (- (deref checking) 10))))
T1 committed
```

1
ı

ref	v0	v1	
checking	100	90	
savings	500	510	

T1: 0

ref	val	write?
checking	90	✓
savings	510	1

T2: 0

ref	val	write?
checking	100	
savings	500	

	_
	-1
	- 1

ref	v0	v1	
checking	100	90	
savings	500	510	

T1: 0

ref	val	write?
checking	90	✓
savings	510	✓

T2: 0

ref	val	write?
checking	90	✓
savings	500	

```
(def checking (ref 100))
(def savings (ref 500))
(future
  (dosync
    (ref-set checking (- (deref checking) 10))
    (ref-set savings (+ (deref savings) 10))))
(future
  (dosync
    (println "You own €" (+ (deref checking)
                             (deref savings)))
    (ref-set checking (- (deref checking) 10))))
commit T2:
```

1. copy written values (if no more recent values exist)

ref	v0	v1	
checking	100	90	
savings	500	510	

D 1

ref	val	write?
checking	90	1
savings	510	1

T1: 0

ref	val	write?
checking	90	✓
savings	500	

```
(def checking (ref 100))
(def savings (ref 500))
(future
  (dosync
    (ref-set checking (- (deref checking) 10))
    (ref-set savings (+ (deref savings) 10))))
(future
  (dosync
    (println "You own €" (+ (deref checking)
                             (deref savings)))
    (ref-set checking (- (deref checking) 10))))
commit T2:
```

1. copy written values (if no more recent values exist)

O 1			T1: 0				T2: 0		
ref	v0	v1	ref	val	write?		ref	val	write
checking	100	90)	checking	90	1	A STATE OF THE STA	checking	90	(V
savings	500	510	savings	510	1		savings	500	

Latest version is more recent than read point of transaction: abort23

	1

ref	v0	v1	
checking	100	90	
savings	500	510	

T1: 0

ref	val	write?
checking	90	1
savings	510	1

T2: **1**

ref	val	write?

1	1
	ı

ref	v0	v1	
checking	100	90	
savings	500	510	

T1: 0

ref	val	write?
checking	90	1
savings	510	1

T2: 1

ref	val	write?
checking	90	
savings	510	

	1
	ı

ref	v0	v1	
checking	100	90	
savings	500	510	

T1: 0

ref	val	write?
checking	90	✓
savings	510	1

T2: 1

ref	val	write?
checking	80	1
savings	510	

```
(def checking (ref 100))
(def savings (ref 500))
(future
  (dosync
    (ref-set checking (- (deref checking) 10))
    (ref-set savings (+ (deref savings) 10))))
(future
  (dosync
    (println "You own €" (+ (deref checking)
                             (deref savings)))
    (ref-set checking (- (deref checking) 10))))
commit T2:
```

1. copy written values (if no more recent values exist)

ref	v0	v1	v2	
checking	100	90	80	
savings	500	510		

1

ref	val	write?
checking	90	1
savings	510	1

T1: 0

ref	val	write?
checking	80	✓
savings	510	

```
(def checking (ref 100))
(def savings (ref 500))
(future
  (dosync
    (ref-set checking (- (deref checking) 10))
    (ref-set savings (+ (deref savings) 10))))
(future
  (dosync
    (println "You own €" (+ (deref checking)
                             (deref savings)))
    (ref-set checking (- (deref checking) 10))))
commit T2:
```

- 1. copy written values (if no more recent values exist)
- 2. increase global write-point

_			
ref	v0	v1	v2
checking	100	90	80
savings	500	510	

P 2

ref val write?

checking 90 ✓

savings 510 ✓

T1: 0

ref val write?

checking 80 ✓

savings 510

```
(def checking (ref 100))
(def savings (ref 500))
(future
  (dosync
    (ref-set checking (- (deref checking) 10))
    (ref-set savings (+ (deref savings) 10))))
(future
  (dosync
    (println "You own €" (+ (deref checking)
                             (deref savings)))
    (ref-set checking (- (deref checking) 10))))
T2 committed
```

7
_

ref	v0	v1	v2	
checking	100	90	80	
savings	500	510		

T1: 0

ref	val	write?
checking	90	1
savings	510	1

T2: 1

ref	val	write?
checking	80	√
savings	510	

Multi-version concurrency control (MVCC)

Multiple versions of refs stored

More parallelism: T2 can read older version even when T1 has updated the same ref

More memory usage ⇒ keep limited history

Optimization: barging detects write-write conflicts early

STM vs locking

```
(def checking (ref 100))
(def savings (ref 500))
(dosync
   (ref-set checking 10)
   (ref-set savings 20))
```

- no deadlocks or livelocks
- no side-effects (retry)
- more overhead
- transactions can be nested
- "optimistic": only take locks when needed

- prone to deadlocks: lock in the right order
- side-effects OK
- less overhead
- nesting can lead to deadlocks
- "pessimistic": all locks at the start

Overview

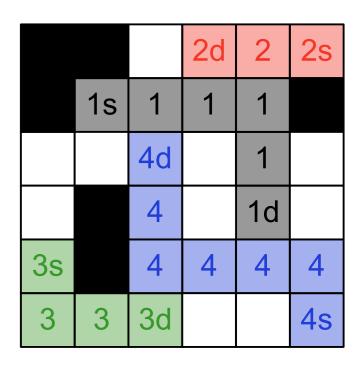
- 1. Crash course: futures
- 2. Crash course: transactions
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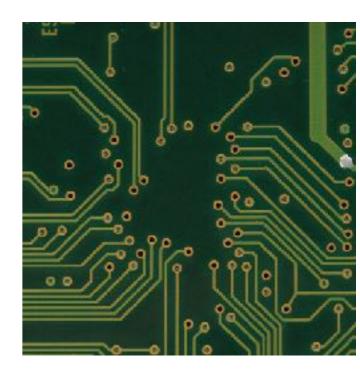
Nesting futures & transactions

```
(future
                             (future
   (future
                                (dosync
     ...))
                                  ...))
Nested task parallelism
                             Transactions
(dosync
                             (dosync
   (future
                                (dosync
     ...))
                                  ...))
                             Nested transactions (open/closed)
In-transaction parallelism
```

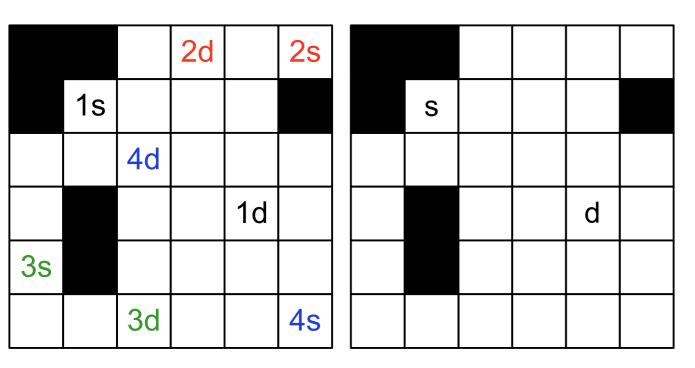
Example: Labyrinth

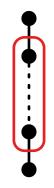
			2d		2s
	1s				
		4d			
				1d	
3s					
		3d			4s





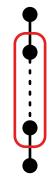
Example: Labyrinth





Example: Labyrinth

			2d		2s			2	3		
	1s						S	1	2	3	
		4d				2	1	2	3		
				1d						d	
3s											
		3d			4s						





		2	3		
	Ø	1	2	3	
2	1	2	3		
				d	

		2	3	4	5
	3	1	2	ى	
2	1	2	3	4	5
3		3	4	d	
4		4	5		
5		5			

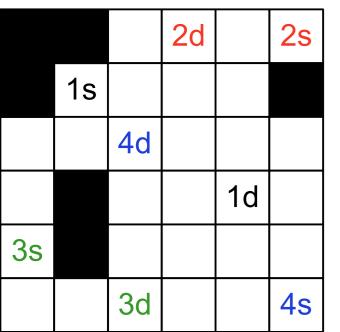
```
    2d
    2s

    1s
    1
    1

    4d
    1

    1d

    3s
    4s
```



		2	3		
	S	1	2	3	
2	1	2	3		
				d	

		2	3	4	5
	8	1	2	رگ	
2	1	2	3	4	5
3		3	4	d	
4		4	5		
5		5			

			2d		2s
	1s	1	1	1	
		4d		1	
				1d	
3s					
		3d			4s

```
(for [[src dst] input-pairs]
  (fetufeocal-grid (copy grid)]
    (depund src dst local-grid)
    (add-path grid
        (traceback local-grid dst))))
```



		2	3		
	S	1	2	3	
2	1	2	3		
				d	

			2	3	4	5
		3	+	CN	P	
	2	1	2	3	4	5
	3		3	4	d	
Ì	4		4	5		
	5		5			

(expand-point (first q) grid)))))))

```
      1s
      1
      1
      1

      4d
      4
      4
      4

      1d
      4

      3s
      4
      4s
```

			2d		2s	
	1s					
		4d				
				1d		
3s						
		3d			4s	

		2	3		
	S	1	2	3	
2	1	2	3		
				d	

		2	3	4	5
	8	1	2	7	
2	1	2	3	4	5
3		3	4	d	
4		4	5		
5		5			

(expand-point (first q) grid)))))))

```
2d 2s

1s 1 1 1

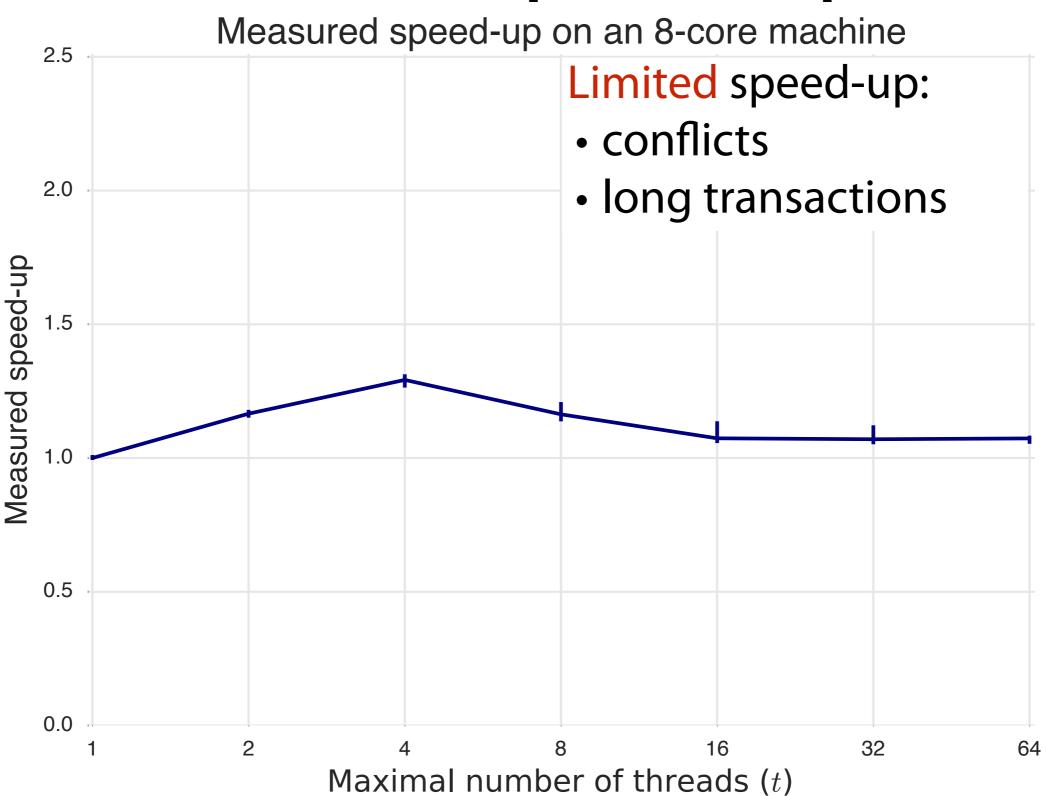
4d 1

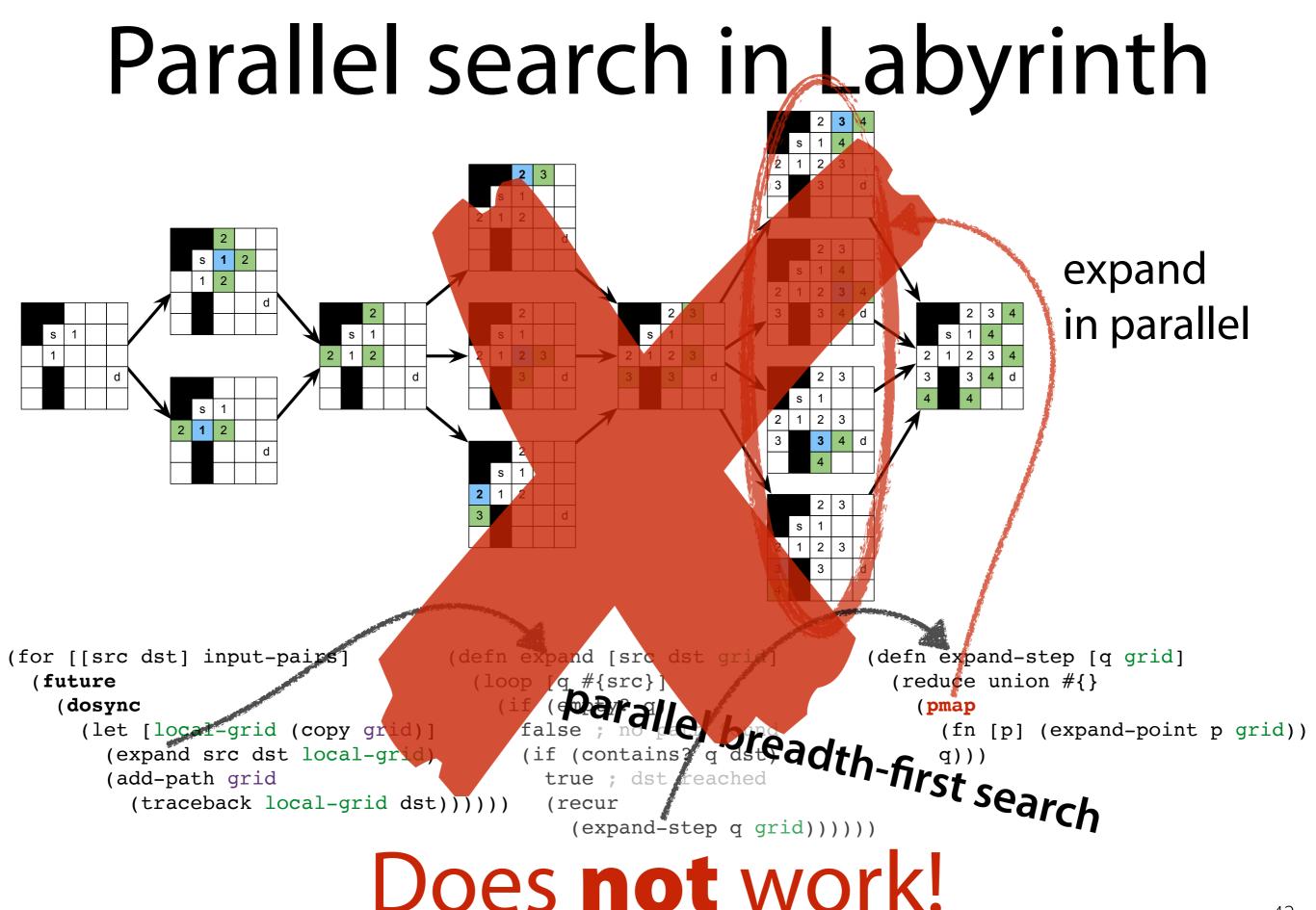
4 1d

3s 4 4 4 4

3d 4s
```

Labyrinth has limited speed-up





Problems when creating threads in a transaction

- Threads in transaction do not share context (Clojure, ScalaSTM)
 - ⇒ no access to transactional state
 - ⇒ serializability violated

```
(dosync
(future
(ref-set ...)))
```

```
(dosync

(future

(dosync

(ref-set ...))))
```

- Threads in transaction prohibited (Haskell)
 - ⇒ parallelism limited

```
atomically $
do { forkIO ... }
```

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Transactional Tasks

Parallelism in transaction

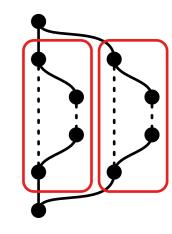
⇒ Transactional task = thread created in transaction

Task can access transactional variables

⇒ Task adopts encapsulating transactional context

Isolation between tasks

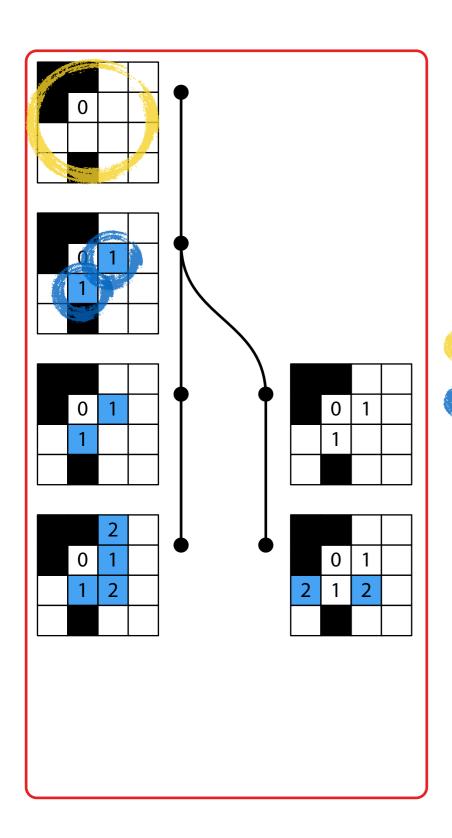
⇒ Tasks work on conceptual copy



Serializability

→ All tasks should join before transaction commits On conflict, all tasks abort

Task = snapshot + store



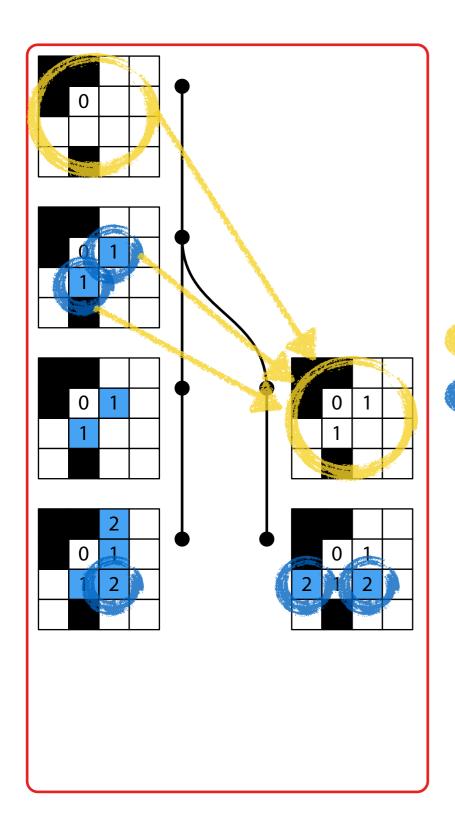
```
(dosync
  (ref-set ... 1)
```

Each transactional task contains:

snapshot: transactional state on creation

Tocal store local modifications

future creates isolated task

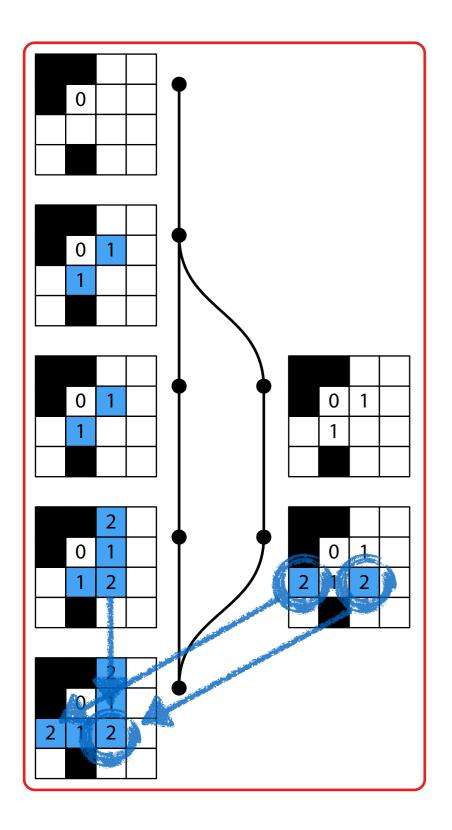


```
(dosync
  (ref-set ... 1)
  (future
        (ref-set ... 2))
  (ref-set ... 2)
```

Each transactional task contains:

snapshot: transactional state on creation local store local modifications

deref merges changes



```
(dosync
...
  (deref child))
```

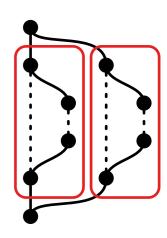
merge local store of child into parent

Conflict resolution function: (ref 0 resolve)

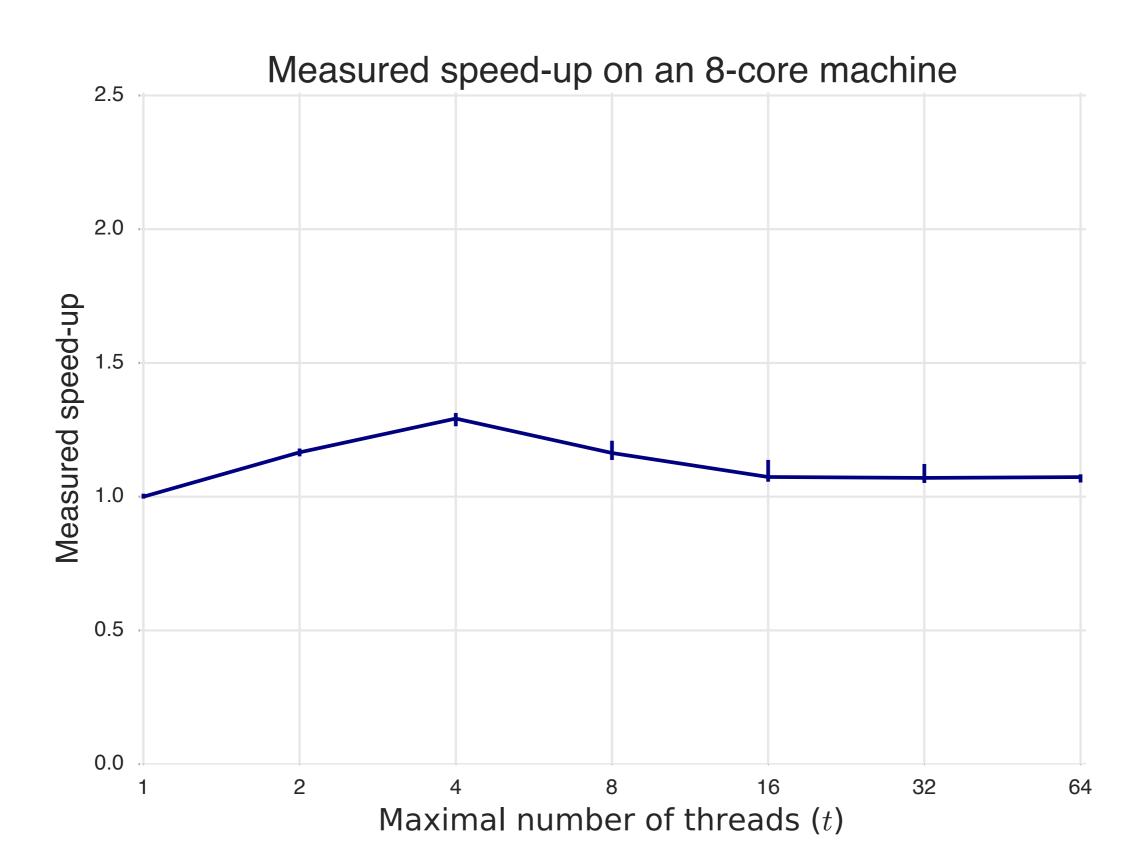
```
(defn resolve [o p c] c)
(defn resolve [o p c] p)
(defn resolve [o p c] (min p c))
(defn resolve [o p c] (+ p c))
(defn resolve [o p c] (error "merge conflict"))
```

Properties of transactional tasks

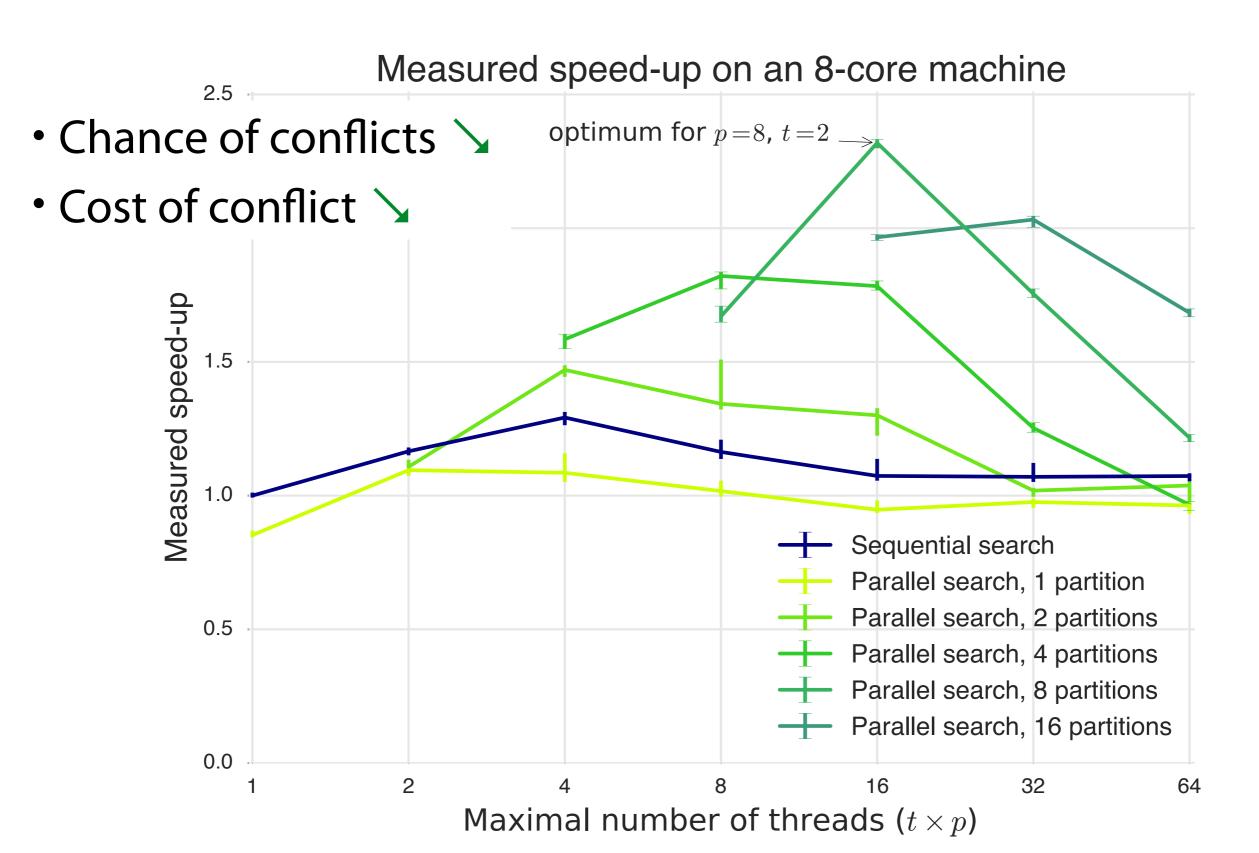
- In-transaction parallelism possible
- Serializability of transactions
- Coordination of tasks: all or none
- In-transaction determinacy



Evaluation: Labyrinth

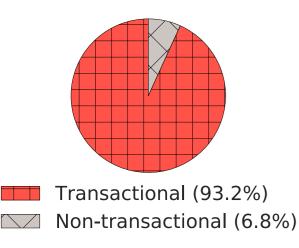


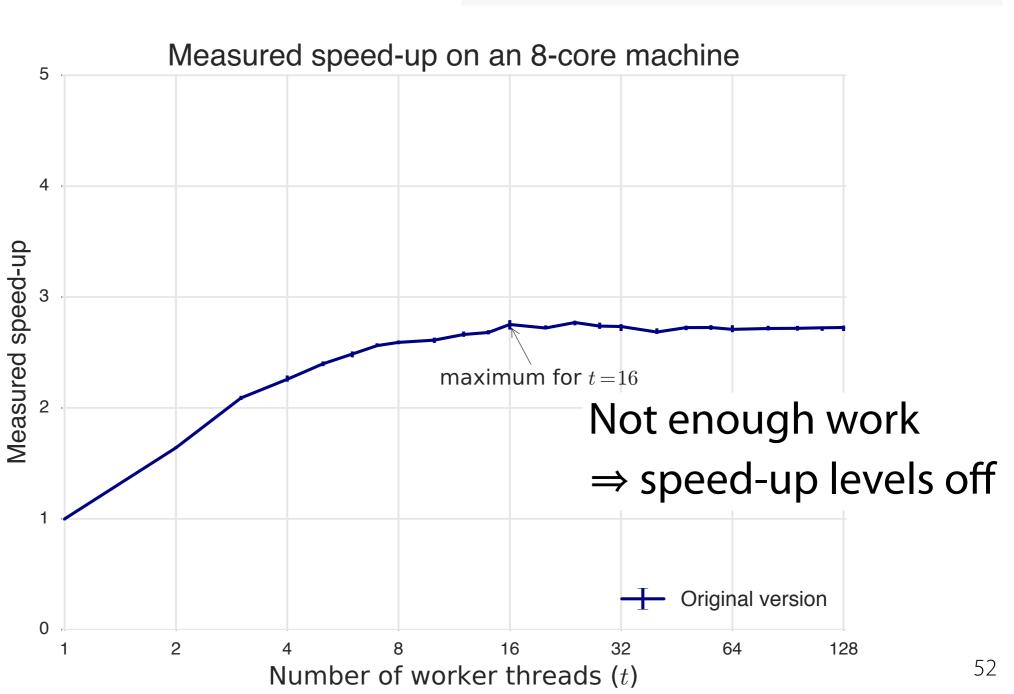
Evaluation: Labyrinth



Evaluation: Bayes

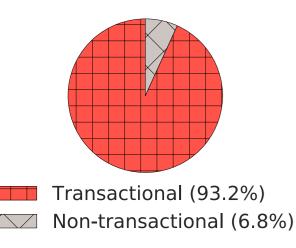
Time spent in transaction (in learning phase)

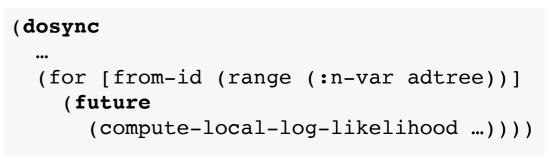


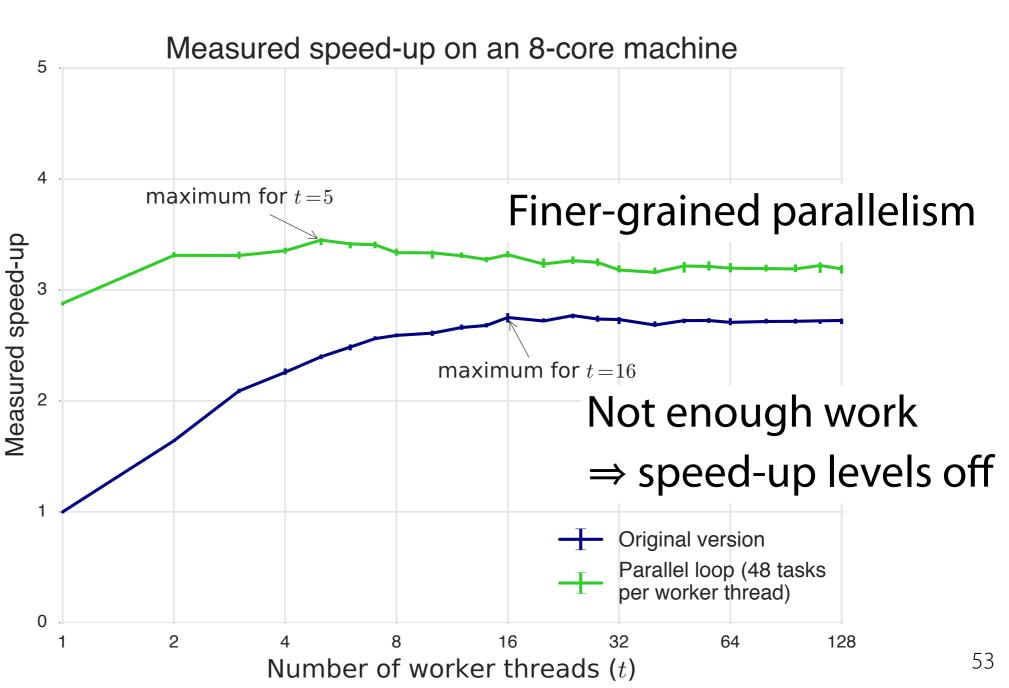


Evaluation: Bayes

Time spent in transaction (in learning phase)







Insights from experiments

- Labyrinth: parallelize search algorithm
 - ⇒ fewer & cheaper conflicts
- Bayes: more fine-grained parallelism
 - ⇒ better exploit hardware
- Low developer effort (re-use existing concepts)
- Suitable for applications with long transactions

Implementation

Fork of Clojure

https://github.com/jswalens/transactional-futures/ http://soft.vub.ac.be/~jswalens/ecoop-2016-artifact/

Summary

Parallelism in a transaction is useful for programs with long transactions But currently:

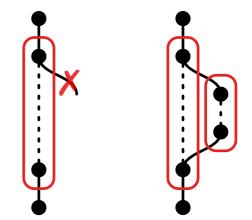
- → not allowed (Haskell)
- ➤ not serializable (Clojure, Scala)

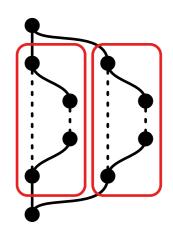
Idea: transactional tasks

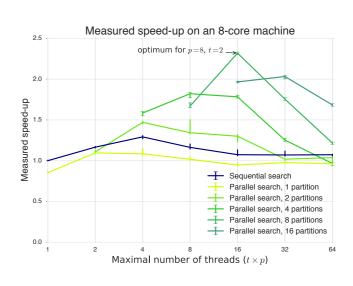
- safe access to encapsulating transaction
- serializable, coordinated, determinate

Benefits:

- finer-grained parallelism ⇒ speed-up
- low developer effort







https://github.com/jswalens/transactional-futures/ http://soft.vub.ac.be/~jswalens/ecoop-2016-artifact/

Transactional Tasks vs. Nested Parallel Transactions

Guarantees in transaction

- NPT: (dosync (future ...)) → race conditions possible
- NPT: (dosync (future (dosync ...))) → serializable, last writer wins (not deterministic)
- TT: conflict resolution → in-transaction determinacy (but may need to define resolution function)

<u>Performance</u>

- NPT: roll back and retry subtransaction
- TT: resolve conflict
- → different performance characteristics depending on application (chance of conflicts between threads in tx)

STAMP

Application	Instructions	Time
	/tx (mean)	in tx
labyrinth	219,571	100%
bayes	60,584	83%
yada	9,795	100%
vacation-high	3,223 ●	86% ●
genome	1,717 ●	<i>97</i> % ●
intruder	$330 \bigcirc$	33% \blacksquare
kmeans-high	$117 \bigcirc$	7%
ssca2	$50 \bigcirc$	17%

Coarse-grained parallelism

between parts of the application

Transactions

- Conflicts span multiple variables
- Difficult to define conflict resolution functions
- Chance of conflicts depends on application
- ⇒ resolve high-level conflicts using serializability

Fine-grained parallelism

within a part of the application

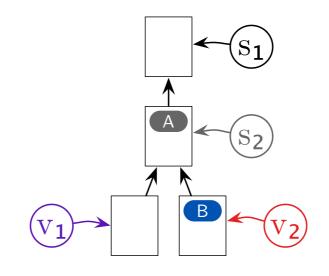
Transactional tasks

- Conflict affects single variable
- Define conflict resolution function based on algorithm
- Conflicts likely, so rollback bad for performance
- ⇒ resolve low-level conflicts using conflict resolution functions

Implementation details

(a) Code example.

(b) Data after step 3.



(c) Data after step 7.

