Transactional Tasks Parallelism in Software Transactions

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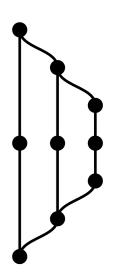
Futures for parallelism

```
(fork e) returns f
(join f) returns result of e
```

```
(defn fib [n]
  (if (< n 2)
        n
        (let [a (fib (- n 1))
             b (fib (- n 2))]
        (+ a b))))</pre>
```

Futures for parallelism

```
(fork e) returns f
(join f) returns result of e
```



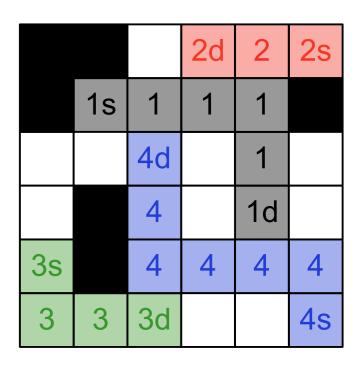
Transactions for shared memory

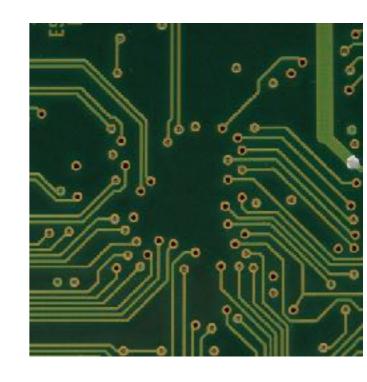
```
serializability
(ref v)
(atomic e)
(deref r)
(ref-set r v)
(def checking (ref 100))
(def savings (ref 500))
(fork
 (atomic
    (ref-set checking (- (deref checking) 10))
    (ref-set savings (+ (deref savings) 10))))
(fork
  (atomic
    (println "You own €" (+ (deref checking)
                            (deref savings)))))
```

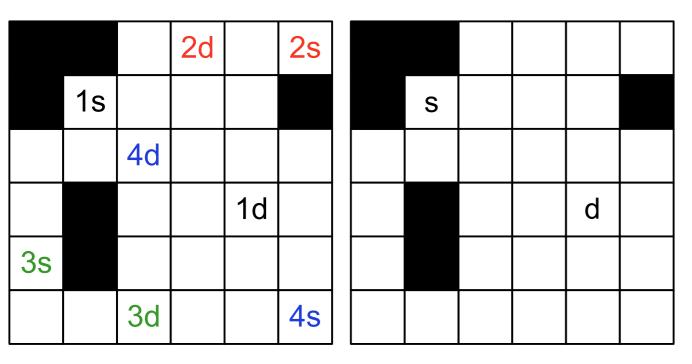
Nesting futures & transactions

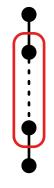
```
(fork
                             (fork
   (fork
                                (atomic
     ...))
                                   ...))
Nested task parallelism
                             Transactions
                             (atomic
(atomic
   (fork
                                (atomic
     ...))
                                   ...))
                             Nested transactions (open/closed)
In-transaction parallelism
```

			2d		2s
	1s				
		4d			
				1d	
3s					
		3d			4s

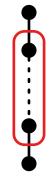


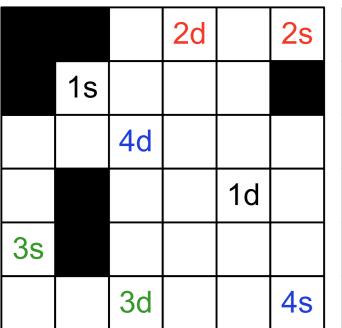






			2d		2s			2	3		
	1s						S	1	2	3	
		4d				2	1	2	3		
				1d						d	
3s											
		3d			4s						

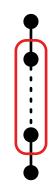


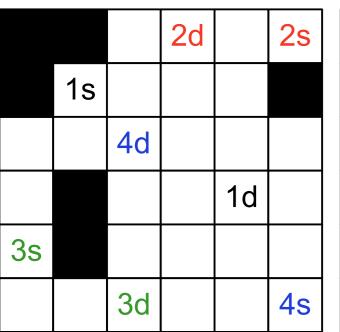


		2	3		
	S	1	2	3	
2	1	2	3		
				d	

		2	3	4	5
	3	1	2	ى	
2	1	2	3	4	5
3		3	4	d	
4		4	5		
5		5			

			2d		2s
	1s	1	1	1	
		4d		1	
				1d	
3s					
		3d			4s





		2	3		
	S	1	2	3	
2	1	2	3		
				d	

		2	3	4	5
	8	1	2	7	
2	1	2	3	4	5
3		3	4	d	
4		4	5		
5		5			

```
    2d
    2s

    1s
    1
    1

    4d
    1

    1d

    3s
    4s
```



		2	3		
	S	1	2	3	
2	1	2	3		
				d	

		2	3	4	5
	3	1	2	ي	
2	1	2	3	4.	5
3		3	4	d	
4		4	5		
5		5			

(expand-point (first q) grid)))))))

```
      2d
      2s

      1s
      1
      1

      4d
      4
      4

      1d
      4

      3s
      4

      3d
      4s
```

			2d		2s	
	1s					
		4d				
				1d		
3s						
		3d			4s	

		2	3		
	S	1	2	3	
2	1	2	3		
				d	

		2	3	4	5
	8	1	2	()	
2	1	2	3	4	5
3		3	4	d	
4		4	5		
5		5			

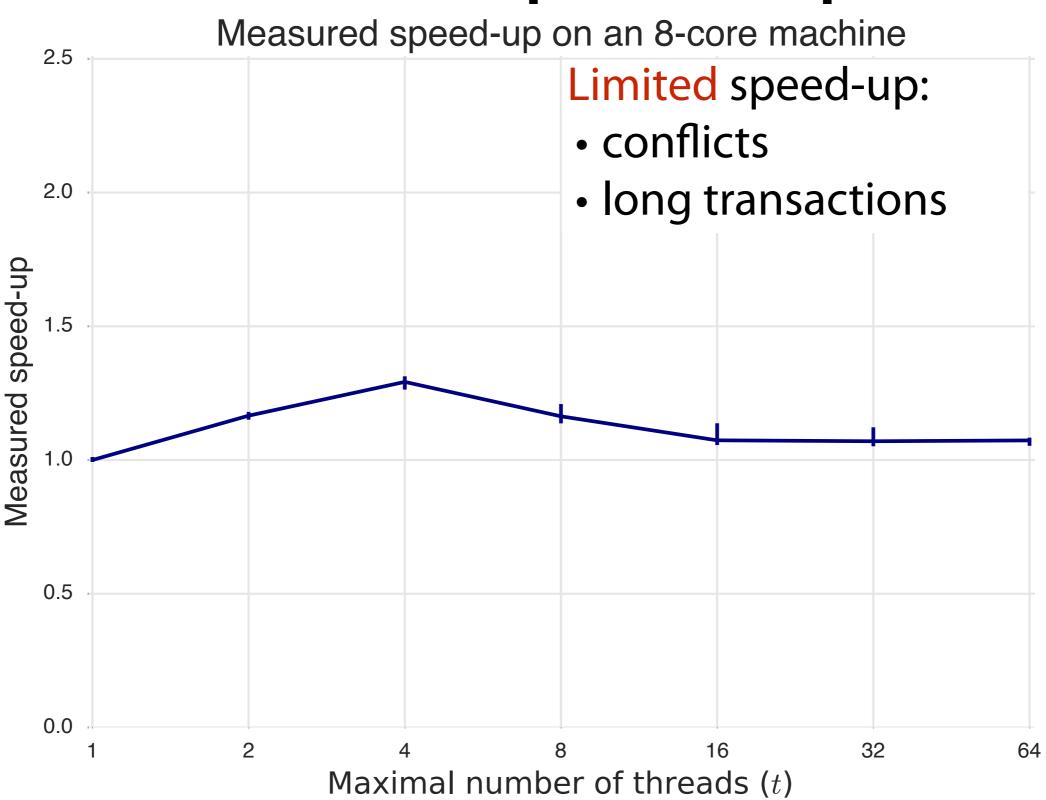
(expand-point (first q) grid))))))

```
2d
                     2s
    1s
        4d
                 1d
3s
        3d
                     4s
```

```
(for [[src dst] input-pairs]
  (fork
     (atomic
       (let [local-grid (copy grid)]
          (expand src dst local-grid)
          (add-path grid
            (traceback local-grid dst))))))
                           (defe expand [src dst grid] (loop [m (list src)]
                                     concat
```

```
defn expand-point [pt grid]
 (atomic
   (let [neighbors ...]
     (for [n neighbors]
       (ref-set n ...))
     neighbors)))
```

Labyrinth has limited speed-up



Parallel search in Labyrinth expand in parallel (for [[src dst] input-pairs (defn expand-step [q grid] (defn expand [src dst g [q #{src} (fork (redu∉e union #{} false; no per adth-first search (atomic (let [local-grid (copy gr (fn [p] (expand-point p grid)) (expand src dst local-grid (add-path grid (traceback local-grid dst))))) (expand-step q grid))))))

Does **not** work!

Problems when creating threads in a transaction

- Threads in transaction do not share context (Clojure, ScalaSTM)
 - ⇒ no access to transactional state
 - ⇒ serializability violated

```
(atomic
  (fork
```

```
(atomic
  (fork
    (atomic (ref-set ...)))
```

- Threads in transaction prohibited (Haskell)
 - ⇒ parallelism limited

```
atomically
```

Transactional Tasks

Parallelism in transaction

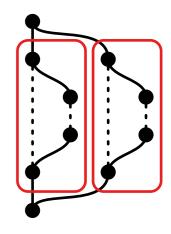
⇒ Transactional task = thread created in transaction

Task can access transactional variables

⇒ Task adopts encapsulating transactional context

Isolation between tasks

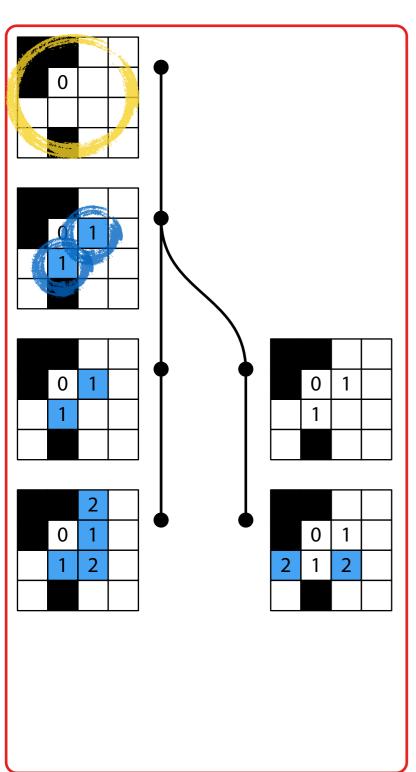
⇒ Tasks work on conceptual copy



Serializability

→ All tasks should join before transaction commits On conflict, all tasks abort

Task = snapshot + store



```
(atomic
  (ref-set ... 1)
```

Each transactional task contains:

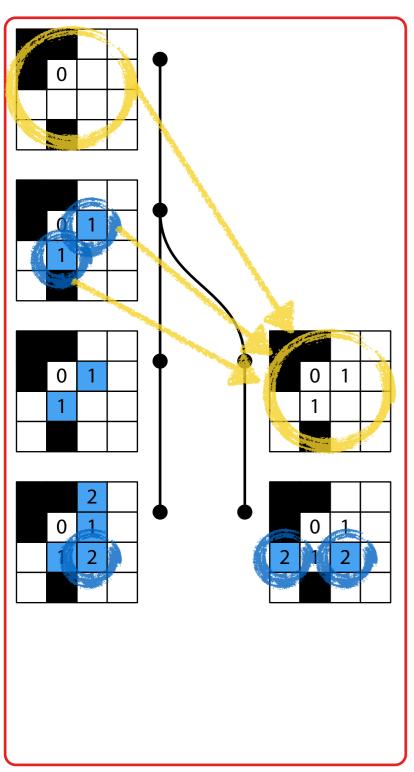
snapshot σ : transactional state on creation

docal store τ local modifications

$$T_{x} \cup \langle f_{p}, \sigma, \tau, F_{s}, F_{j}, \mathcal{E}[fork e] \rangle$$

$$\Rightarrow_{tf} T_{x} \cup \langle f_{p}, \sigma, \tau, F_{s} \cup \{f'\}, F_{j}, \mathcal{E}[f_{c}] \rangle \cup \langle f_{c}, \sigma :: \tau, \varnothing, \varnothing, F_{j}, e \rangle$$
with f_{c} fresh

fork creates isolated task



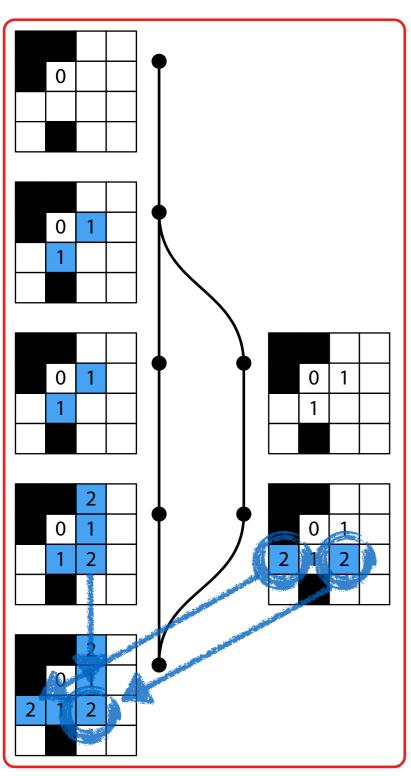
```
(atomic
  (ref-set ... 1)
  (fork
        (ref-set ... 2))
  (ref-set ... 2)
```

Each transactional task contains:

 σ transactional state on creation σ local modifications

```
T_{x} \cup \langle f_{p}, \sigma, \tau, F_{s}, F_{j}, \mathcal{E}[fork e] \rangle
\Rightarrow_{tf} T_{x} \cup \langle f_{p}, \sigma, \tau, F_{s} \cup \{f'\}, F_{j}, \mathcal{E}[f_{c}] \rightarrow \cup \langle f_{c}, \sigma :: \tau | \varnothing \varnothing, F_{j}, e \rangle
with f_{c} fresh
```

join merges changes



```
(atomic
...
  (join child))
```

merge local store τ' of child into parent

```
T_{\mathsf{x}} \cup \langle f_p, \sigma(\tau) \mid F_{\mathsf{s}}, F_{\mathsf{j}}, \qquad \mathcal{E}[\mathsf{join}\ f_c] \rangle \cup \langle f_c, \sigma', \tau', F_{\mathsf{s}}', F_{\mathsf{j}}', v \rangle \ \Rightarrow_{\mathsf{tf}} T_{\mathsf{x}} \cup \langle f_p, \sigma(\tau :: \tau') \mid F_{\mathsf{s}}', F_{\mathsf{j}}' \cup F_{\mathsf{j}}' \cup \{f_c\}, \mathcal{E}[v] \qquad \rangle \cup \langle f_c, \sigma', \tau', F_{\mathsf{s}}', F_{\mathsf{j}}', v \rangle \ \text{if}\ f_c \notin F_{\mathsf{j}} \ \text{and}\ F_{\mathsf{s}}' \subseteq F_{\mathsf{j}}'
```

Conflict resolution function:

(ref 0 resolve)

```
resolve::T×T×T→T

(defn resolve [o p c] c)

(defn resolve [o p c] p)

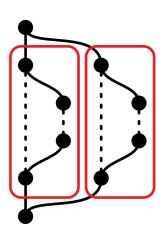
(defn resolve [o p c] (min p c))

(defn resolve [o p c] (+ p c))

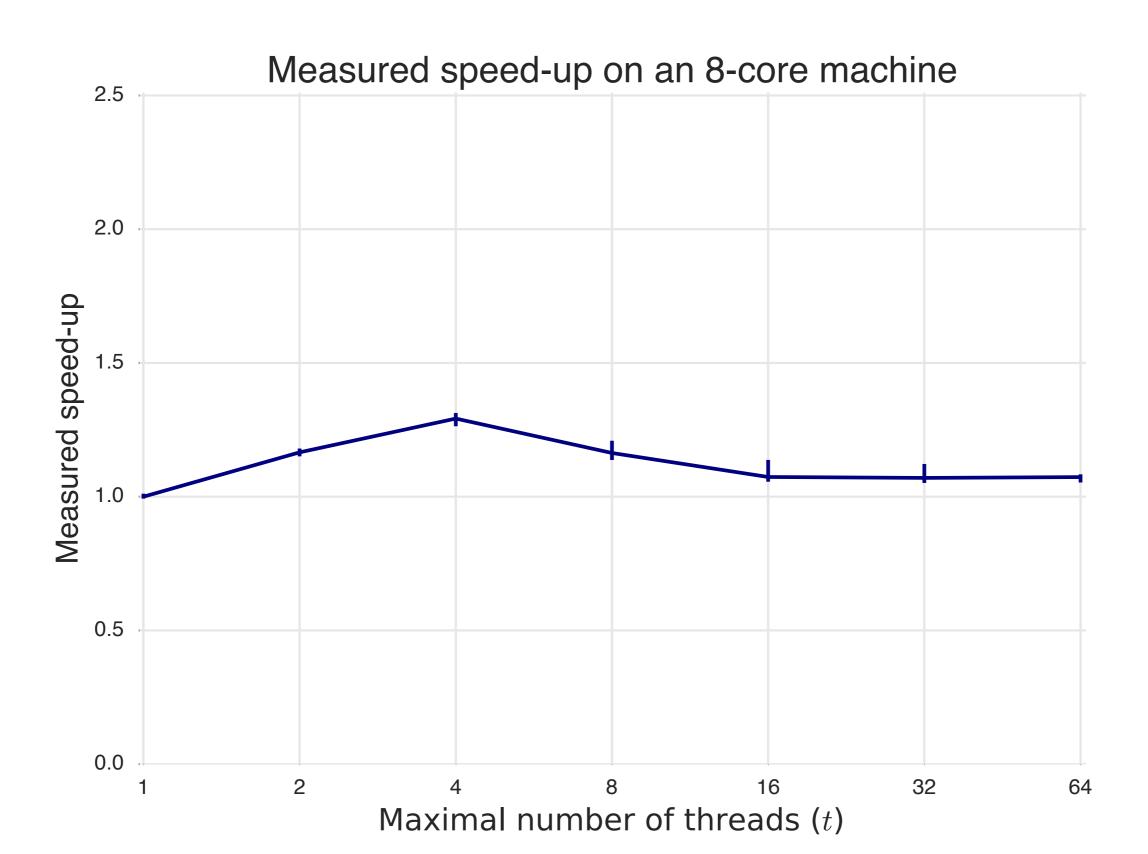
(defn resolve [o p c] (error "merge conflict"))
```

Properties of transactional tasks

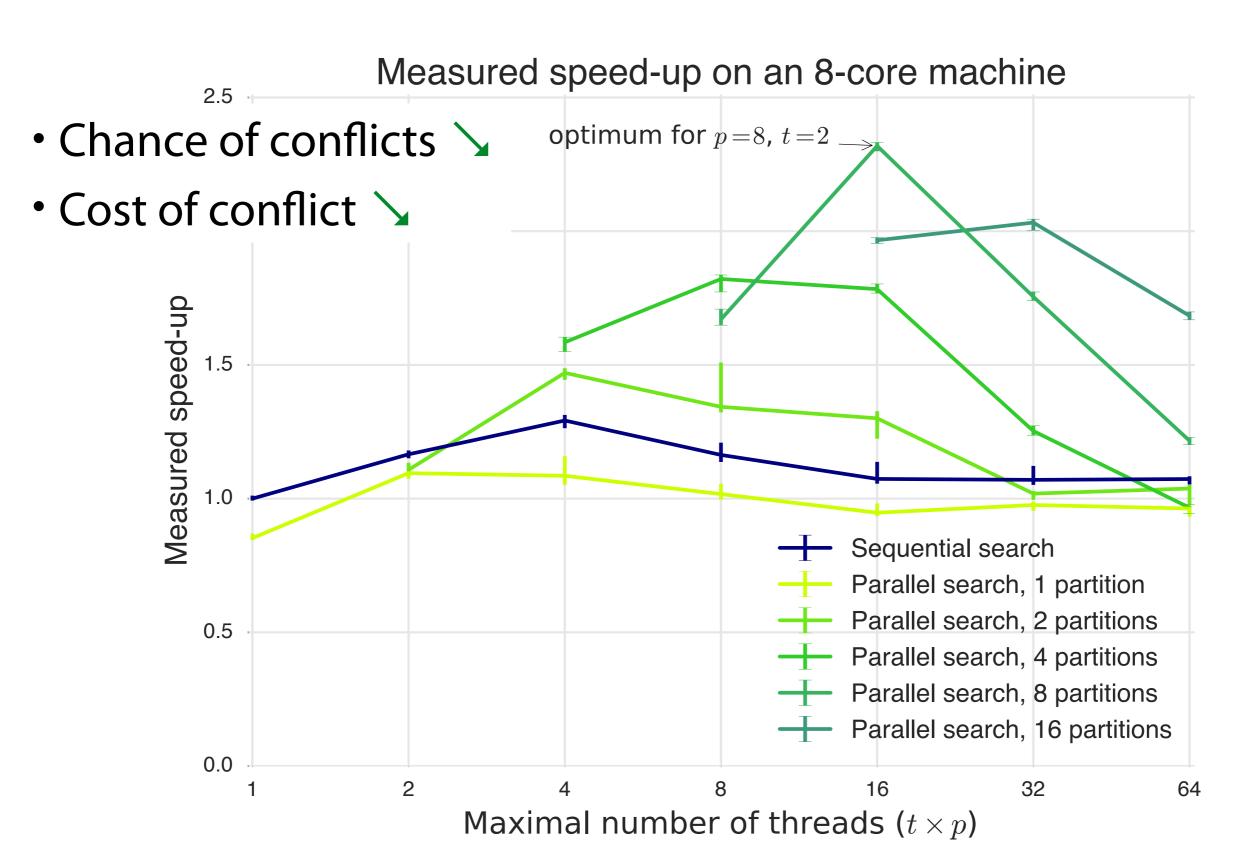
- In-transaction parallelism possible
- Serializability of transactions
- Coordination of tasks: all or none
- In-transaction determinacy



Evaluation: Labyrinth

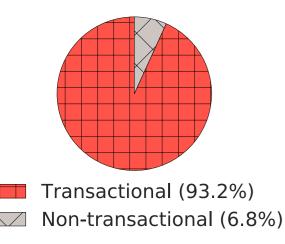


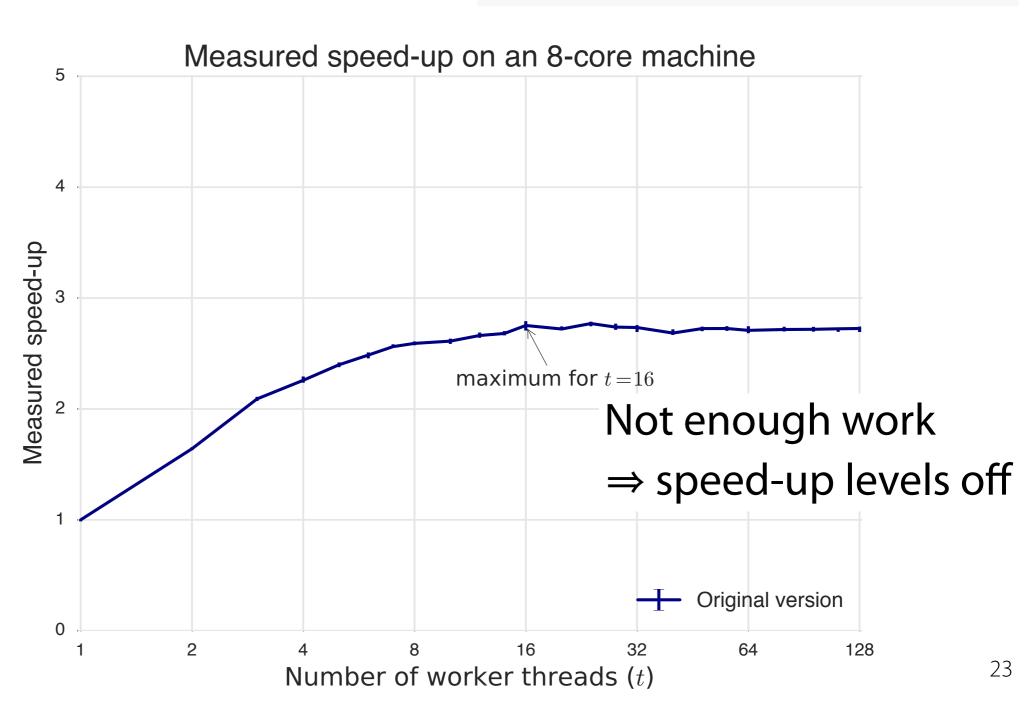
Evaluation: Labyrinth



Evaluation: Bayes

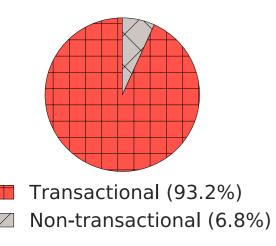
Time spent in transaction (in learning phase)

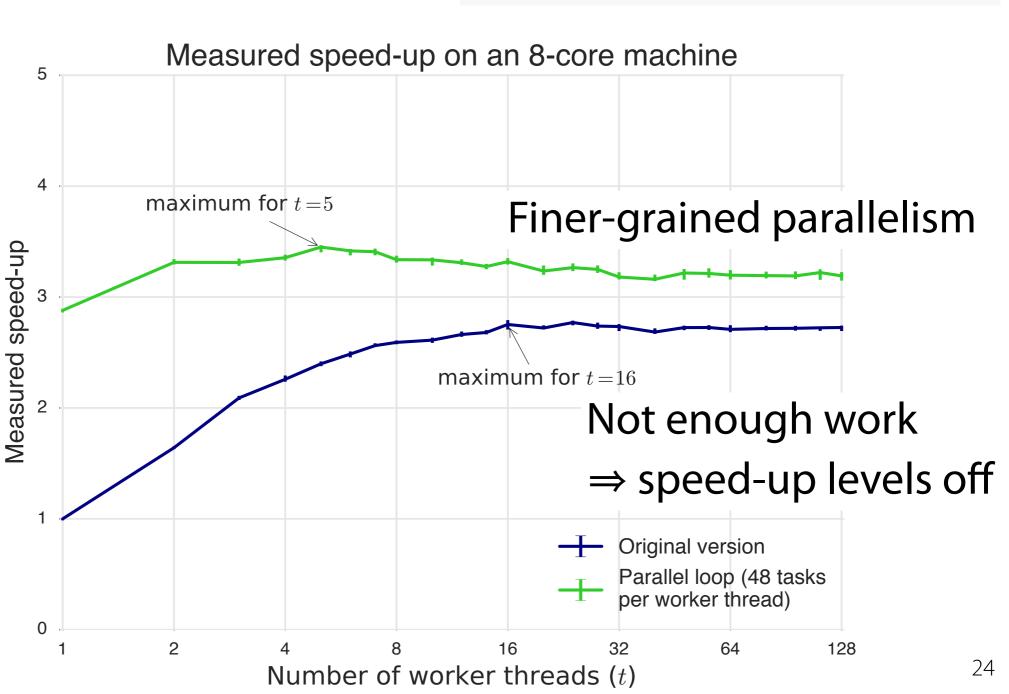




Evaluation: Bayes

Time spent in transaction (in learning phase)





Insights from experiments

- Labyrinth: parallelize search algorithm
 - ⇒ fewer & cheaper conflicts
- Bayes: more fine-grained parallelism
 - ⇒ better exploit hardware
- Low developer effort (re-use existing concepts)
- Suitable for applications with long transactions

Implementation

Fork of Clojure

https://github.com/jswalens/transactional-futures/

http://soft.vub.ac.be/~jswalens/ecoop-2016-artifact/

Details in paper



Summary

Parallelism in a transaction is useful for programs with long transactions
But currently:

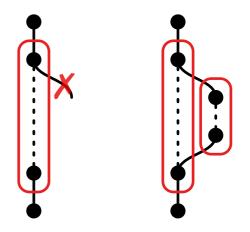
- ∧ not allowed (Haskell)
- ➤ not serializable (Clojure, Scala)

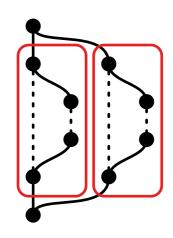
Idea: transactional tasks

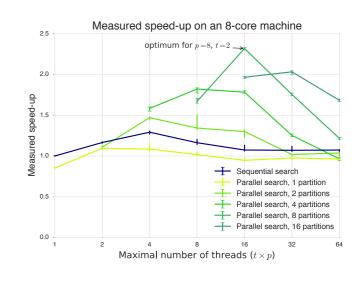
- safe access to encapsulating transaction
- serializable, coordinated, determinate

Benefits:

- finer-grained parallelism ⇒ speed-up
- low developer effort







Transactional Tasks vs. Nested Parallel Transactions

Guarantees in transaction

- NPT: (atomic (fork ...)) → race conditions possible
- NPT: (atomic (fork (atomic ...))) → serializable, last writer wins (not deterministic)
- TT: conflict resolution → in-transaction determinacy (but may need to define resolution function)

Performance

- NPT: roll back and retry subtransaction
- TT: resolve conflict
- → different performance characteristics depending on application (chance of conflicts between threads in tx)

STAMP

Application	Instructions	Time
	/tx (mean)	in tx
labyrinth	219,571	100%
bayes	60,584	83%
yada	9,795	100%
vacation-high	3,223 ●	86% ●
genome	1,717 ●	97%
intruder	$330 \bigcirc$	33% ●
kmeans-high	$117 \bigcirc$	7% \bigcirc
ssca2	$50 \bigcirc$	17%

Coarse-grained parallelism

between parts of the application

Transactions

- Conflicts span multiple variables
- Difficult to define conflict resolution functions
- Chance of conflicts depends on application
- ⇒ resolve high-level conflicts using serializability

Fine-grained parallelism

within a part of the application

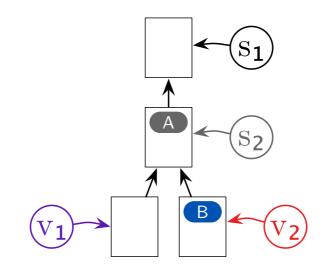
Transactional tasks

- Conflict affects single variable
- Define conflict resolution function based on algorithm
- Conflicts likely, so rollback bad for performance
- ⇒ resolve low-level conflicts using conflict resolution functions

Implementation details

(a) Code example.

(b) Data after step 3.



(c) Data after step 7.

